

var, x, y	term variable		
t	::=		term
		x	variable
		$\lambda x.t$	bind x in t lambda
		$t t'$	app
		if t then t' else t''	conditional
		true	true
		false	false
		(t)	S
		$[t/x]t'$	M
v	::=		value
		$\lambda x.t$	lambda
typ, T	::=		types
		Bool	bool type
		$T_1 \rightarrow T_2$	function types
ctx, Γ	::=		typing context
		\bullet	empty context
		$\Gamma, x : T$	assumption
$terminals$::=		
		λ	
		\longrightarrow	
		\rightarrow	
		\in	
		\neq	
		\equiv_α	
		\equiv_β	
		FV	
		\notin	
		dom	
		\vdash	
		true	
		false	
$formula$::=		
		<i>judgement</i>	
		$x \neq x'$	M
		$x \notin \text{FV}(t)$	M
		$x : T \in \Gamma$	M
		$x \notin \text{dom}(\Gamma)$	M
red	::=		
		$t_1 \longrightarrow t_2$	t_1 reduces to t_2
fv	::=		
		$x \in \text{FV}(t)$	free variable

<i>aeq</i>	::=	$t \equiv_{\alpha} t'$	alpha equivalence
<i>beq</i>	::=	$t \equiv_{\beta} t'$	beta equivalence
<i>typing</i>	::=	$\Gamma \vdash t : T$	Typing rules
<i>judgement</i>	::=	<i>red</i>	
		<i>fv</i>	
		<i>aeq</i>	
		<i>beq</i>	
		<i>typing</i>	
<i>user_syntax</i>	::=	<i>var</i>	
		<i>t</i>	
		<i>v</i>	
		<i>typ</i>	
		<i>ctx</i>	
		<i>terminals</i>	
		<i>formula</i>	

$t_1 \longrightarrow t_2$ t_1 reduces to t_2

$\frac{}{(\lambda x. t_{12}) v_2 \longrightarrow [v_2/x] t_{12}}$	RED_AX_APP
$\frac{t_1 \longrightarrow t'_1}{t_1 t \longrightarrow t'_1 t}$	RED_CTX_APP_FUN
$\frac{t_1 \longrightarrow t'_1}{v t_1 \longrightarrow v t'_1}$	RED_CTX_APP_ARG
$\frac{}{\mathbf{if\ true\ then\ } t_1 \mathbf{\ else\ } t_2 \longrightarrow t_1}$	RED_IF_TRUE
$\frac{}{\mathbf{if\ false\ then\ } t_1 \mathbf{\ else\ } t_2 \longrightarrow t_2}$	RED_IF_FALSE
$\frac{t_1 \longrightarrow t'_1}{\mathbf{if\ } t_1 \mathbf{\ then\ } t_2 \mathbf{\ else\ } t_3 \longrightarrow \mathbf{if\ } t'_1 \mathbf{\ then\ } t_2 \mathbf{\ else\ } t_3}$	RED_IF

$x \in \text{FV}(t)$ free variable

$\frac{}{x \in \text{FV}(x)}$	FV_VAR
$\frac{x \in \text{FV}(t_1)}{x \in \text{FV}(t_1 t_2)}$	FV_APP_L
$\frac{x \in \text{FV}(t_2)}{x \in \text{FV}(t_1 t_2)}$	FV_APP_R

$$\frac{x \in \text{FV}(t) \quad x \neq y}{x \in \text{FV}(\lambda y.t)} \quad \text{FV_LAM}$$

$t \equiv_\alpha t'$ alpha equivalence

$$\frac{}{t \equiv_\alpha t} \quad \text{AEQ_ID}$$

$$\frac{t \equiv_\alpha t'}{t' \equiv_\alpha t} \quad \text{AEQ_SYM}$$

$$\frac{t \equiv_\alpha t' \quad t' \equiv_\alpha t''}{t \equiv_\alpha t''} \quad \text{AEQ_TRANS}$$

$$\frac{t_1 \equiv_\alpha t'_1 \quad t_2 \equiv_\alpha t'_2}{t_1 t_2 \equiv_\alpha t'_1 t'_2} \quad \text{AEQ_APP}$$

$$\frac{t \equiv_\alpha t'}{\lambda x.t \equiv_\alpha \lambda x.t'} \quad \text{AEQ_LAM}$$

$$\frac{x' \notin \text{FV}(t)}{\lambda x.t \equiv_\alpha \lambda x'. [x'/x]t} \quad \text{AEQ_SUBST}$$

$t \equiv_\beta t'$ beta equivalence

$$\frac{}{t \equiv_\beta t} \quad \text{BEQ_ID}$$

$$\frac{t \equiv_\beta t'}{t' \equiv_\beta t} \quad \text{BEQ_SYM}$$

$$\frac{t \equiv_\beta t' \quad t' \equiv_\beta t''}{t \equiv_\beta t''} \quad \text{BEQ_TRANS}$$

$$\frac{t_1 \equiv_\beta t'_1 \quad t_2 \equiv_\beta t'_2}{t_1 t_2 \equiv_\beta t'_1 t'_2} \quad \text{BEQ_APP}$$

$$\frac{t \equiv_\beta t'}{\lambda x.t \equiv_\beta \lambda x.t'} \quad \text{BEQ_LAM}$$

$$\frac{}{(\lambda x.t) t' \equiv_\beta [t'/x]t} \quad \text{BEQ_SUBST}$$

$\Gamma \vdash t : T$ Typing rules

$$\frac{x : T \in \Gamma}{\Gamma \vdash x : T} \quad \text{TYPING_VAR}$$

$$\frac{\Gamma, x : T_1 \vdash t : T_2}{\Gamma \vdash \lambda x.t : T_1 \rightarrow T_2} \quad \text{TYPING_ABS}$$

$$\frac{\Gamma \vdash t_1 : T_1 \rightarrow T_2 \quad \Gamma \vdash t_2 : T_1}{\Gamma \vdash t_1 t_2 : T_2} \quad \text{TYPING_APP}$$

$$\frac{}{\Gamma \vdash \text{true} : \text{Bool}} \quad \text{TYPING_TRUE}$$

$$\frac{}{\Gamma \vdash \text{false} : \text{Bool}} \quad \text{TYPING_FALSE}$$

$$\frac{\Gamma \vdash t_1 : \text{Bool} \quad \Gamma \vdash t_2 : T_1 \quad \Gamma \vdash t_3 : T_1}{\Gamma \vdash \text{if } t_1 \text{ then } t_2 \text{ else } t_3 : T_1} \quad \text{TYPING_IF}$$

Definition rules: 28 good 0 bad
 Definition rule clauses: 55 good 0 bad